# Chapter 6 Adversarial Search

CS 461 – Artificial Intelligence Pinar Duygulu Bilkent University, Spring 2008

Slides are mostly adapted from AIMA and MIT Open Courseware

### Outline

- Games
- Optimal decisions
- Minimax algorithm
- $\alpha$ - $\beta$  pruning
- Imperfect, real-time decisions

2

### Games

- Multi agent environments : any given agent will need to consider the actions of other agents and how they affect its own welfare.
- The unpredictability of these other agents can introduce many possible contingencies
- There could be competitive or cooperative environments
- Competitive environments, in which the agent's goals are in conflict require adversarial search – these problems are called as games

- In game theory (economics), any multiagent environment (either cooperative or competitive) is a game provided that the impact of each agent on the other is significant
- AI games are a specialized kind deterministic, turn taking, two-player, zero sum games of perfect information
- In our terminology deterministic, fully observable environments with two agents whose actions alternate and the utility values at the end of the game are always equal and opposite (+1 and -1)

### Games – history of chess playing

- 1949 Shannon paper originated the ideas
- 1951 Turing paper hand simulation
- 1958 Bernstein program
- 1955-1960 Simon-Newell program
- 1961 Soviet program
- 1966 1967 MacHack 6 defeated a good player
- 1970s NW chess 4.5
- 1980s Cray Bitz
- 1990s Belle, Hitech, Deep Thought,
- 1997 Deep Blue defeated Garry Kasparov

- Initial state: initial board position and player
- Operators: one for each legal move
- Goal states: winning board positions
- Scoring function: assigns numeric value to states
- Game tree: encodes all possible games
- We are not looking for a path, only the next move to make (that hopefully leads to a winning position)
- Our best move depends on what the other player does

### Partial Game Tree for Tic-Tac-Toe



### Game tree (2-player, deterministic, turns)



- In a normal search problem, the optimal solution would be a sequence of moves leading to a goal state a terminal state that is a win
- In a game, MIN has something to say about it and therefore MAX must find a contingent strategy, which specifies
  - MAX's move in the initial state,
  - then MAX's moves in the states resulting from every possible response by MIN,
  - then MAX's moves in the states resulting from every possible response by MIN to those moves

- ...

• An optimal strategy leads to outcomes at least as good as any other strategy when one is playing an infallible opponent

### Minimax

- Perfect play for deterministic games
- Idea: choose move to position with highest minimax value = best achievable payoff against best play
- E.g., 2-ply game:



- Given a game tree, the optimal strategy can be determined by examining the minimax value of each node (MINIMAX-VALUE(n))
- The minimax value of a node is the utility of being in the corresponding state, assuming that both players play optimally from there to the end of the game
- Given a choice, MAX prefer to move to a state of maximum value, whereas MIN prefers a state of minimum value

11

function MINIMAX-DECISION(state) returns an action

```
v \leftarrow \text{MAX-VALUE}(state)
return the action in SUCCESSORS(state) with value v
```

function MAX-VALUE(state) returns a utility value

if TERMINAL-TEST(*state*) then return UTILITY(*state*)

```
v \leftarrow -\infty
```

for a, s in SUCCESSORS(state) do

```
v \leftarrow Max(v, MIN-VALUE(s))
```

return v

function MIN-VALUE(state) returns a utility value

if TERMINAL-TEST(*state*) then return UTILITY(*state*)

```
v \leftarrow \infty
```

```
for a, s in SUCCESSORS(state) do
```

```
v \leftarrow MIN(v, MAX-VALUE(s))
```

```
return v
```

### Minimax

### MINIMAX-VALUE(root) = $\max(\min(3,12,8), \min(2,4,6), \min(14,5,2))$ = $\max(3,2,2)$ = 3



### Properties of minimax

- <u>Complete?</u> Yes (if tree is finite)
- <u>Optimal?</u> Yes (against an optimal opponent)
- <u>Time complexity?</u> O(b<sup>m</sup>)
- <u>Space complexity?</u> O(bm) (depth-first exploration)
- For chess, b ≈ 35, m ≈100 for "reasonable" games
   → exact solution completely infeasible

### Tree Player and Non-zero sum games



• It is possible to compute the correct minimax decision without looking at every node in the game tree













### Properties of $\alpha$ - $\beta$

- Pruning does not affect final result
- Good move ordering improves effectiveness of pruning
- With "perfect ordering," time complexity = O(b<sup>m/2</sup>)
   → doubles depth of search
- A simple example of the value of reasoning about which computations are relevant (a form of metareasoning)

# Why is it called $\alpha$ - $\beta$ ?

- α is the value of the best (i.e., highest-value) choice found so far at any choice point along the path for max
- If v is worse than α, max will avoid it
   → prune that branch
- Define β similarly for *min*



function ALPHA-BETA-SEARCH(state) returns an action inputs: state, current state in game

```
v \leftarrow \text{MAX-VALUE}(state, -\infty, +\infty)
return the action in SUCCESSORS(state) with value v
```

function MAX-VALUE(*state*,  $\alpha$ ,  $\beta$ ) returns a utility value inputs: *state*, current state in game

 $lpha_{\text{r}}$  the value of the best alternative for  $_{ ext{MAX}}$  along the path to state

eta, the value of the best alternative for  $_{
m MIN}$  along the path to state

```
if TERMINAL-TEST(state) then return UTILITY(state)
```

```
v \leftarrow -\infty
for a, s in SUCCESSORS(state) do
v \leftarrow MAX(v, MIN-VALUE(s, \alpha, \beta))
if v \ge \beta then return v
\alpha \leftarrow MAX(\alpha, v)
```

return v

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function MIN-VALUE(state, \alpha, \beta) returns a utility value
inputs: state, current state in game
\alpha, the value of the best alternative for MAX along the path to state
\beta, the value of the best alternative for MIN along the path to state
if TERMINAL-TEST(state) then return UTILITY(state)
v \leftarrow +\infty
for a, s in SUCCESSORS(state) do
v \leftarrow MIN(v, MAX-VALUE(s, \alpha, \beta))
if v \leq \alpha then return v
\beta \leftarrow MIN(\beta, v)
return v
```

### The $\alpha$ - $\beta$ algorithm

### α - β

```
// \alpha = best score for MAX, \beta = best score for MIN
// initial call is MAX-VALUE(state,-\infty, \infty, MAX-DEPTH)
```

```
function MAX-VALUE (state, \alpha, \beta, depth)

if (depth == 0) then return EVAL (state)

for each s in SUCCESSORS (state) do

\alpha = MAX (\alpha, MIN-VALUE (s, \alpha, \beta, depth-1))

if \alpha \ge \beta then return \alpha // cutoff

end

return \alpha
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## $\alpha$ - $\beta$ pruning



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### Alpha-beta Pruning



### Move generation



# Suppose we have 100 secs, explore $10^4$ nodes/sec $\rightarrow 10^6$ nodes per move

Standard approach:

• cutoff test:

e.g., depth limit (perhaps add quiescence search)

- evaluation function
  - = estimated desirability of position

### Evaluation function



### Min-Max



- A typical evaluation function is a linear function in which some set of coefficients is used to weight a number of "features" of the board position.
- For chess, typically linear weighted sum of features

 $Eval(s) = w_1 f_1(s) + w_2 f_2(s) + \ldots + w_n f_n(s)$ 

• e.g.,  $w_1 = 9$  with

 $f_1(s) = (number of white queens) - (number of black queens), etc.$ 

### Evaluation function

S	=	<b>c</b> <sub>1</sub>	x	material
a a a a a a a a a a a a a a a a a a a		<b>c</b> <sub>2</sub>	x	pawn structure
Ŧ		c <sub>3</sub>	x	mobilit <b>y</b>
1		<b>c</b> <sub>4</sub>	x	king safety
2 -		$\mathbf{c}_5$	x	center control
÷				



- "material", : some measure of which pieces one has on the board.
- A typical weighting for each type of chess piece is shown
- Other types of features try to encode something about the distribution of the pieces on the board.

50

*MinimaxCutoff* is identical to *MinimaxValue* except

- *1. Terminal?* is replaced by *Cutoff?*
- 2. Utility is replaced by Eval

```
Does it work in practice?
b^m = 10^6, b=35 \rightarrow m=4
```

4-ply lookahead is a hopeless chess player!

- 4-ply  $\approx$  human novice
- 8-ply  $\approx$  typical PC, human master
- 12-ply  $\approx$  Deep Blue, Kasparov



• The key idea is that the more lookahead we can do, that is, the deeper in the tree we can look, the better our evaluation of a position will be, even with a simple evaluation function. In some sense, if we could look all the way to the end of the game, all we would need is an evaluation function that was 1 when we won and -1 when the opponent won.

- it seems to suggest that brute-force search is all that matters.
- And Deep Blue is brute indeed... It had 256 specialized chess processors coupled into a 32 node supercomputer. It examined around 30 billion moves per minute. The typical search depth was 13ply, but in some dynamic situations it could go as deep as 30.

### Practical issues

### Variable branching



#### Iterative deepening

- order best move from last search first
- use previous backed up value to initialize  $[\alpha, \beta]$
- keep track of repeated positions (transposition tables)

### Horizon effect

- 🗕 quiescence
  - Pushing the inevitable over search horizon

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- Backgammon
  - Involves randomness dice rolls
  - Machine-learning based player was able to draw the world champion human player.
- Bridge
  - Involves hidden information other players' cards and communication during bidding.
  - Computer players play well but do not bid well
- Go
  - No new elements but huge branching factor
  - No good computer players exist

### Deterministic games in practice

- Checkers: Chinook ended 40-year-reign of human world champion Marion Tinsley in 1994. Used a precomputed endgame database defining perfect play for all positions involving 8 or fewer pieces on the board, a total of 444 billion positions.
- Chess: Deep Blue defeated human world champion Garry Kasparov in a sixgame match in 1997. Deep Blue searches 200 million positions per second, uses very sophisticated evaluation, and undisclosed methods for extending some lines of search up to 40 ply.
- Othello: human champions refuse to compete against computers, who are too good.
- Go: human champions refuse to compete against computers, who are too bad. In go, b > 300, so most programs use pattern knowledge bases to suggest plausible moves.

- Games are fun to work on!
- They illustrate several important points about AI
- perfection is unattainable  $\rightarrow$  must approximate
- good idea to think about what to think about