

Priority Queues - 2

CS 202 – Fundamental Structures of
Computer Science II

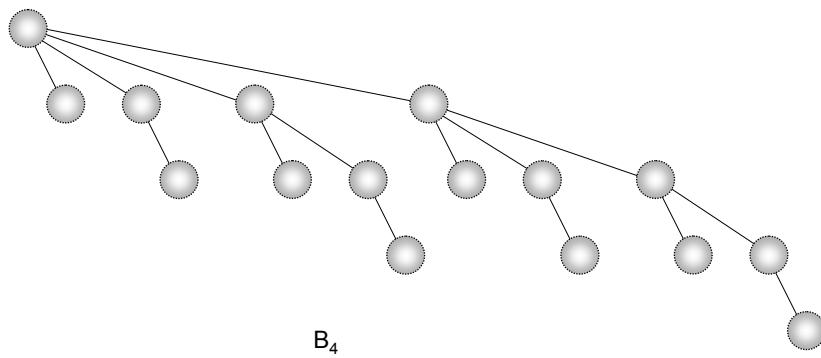
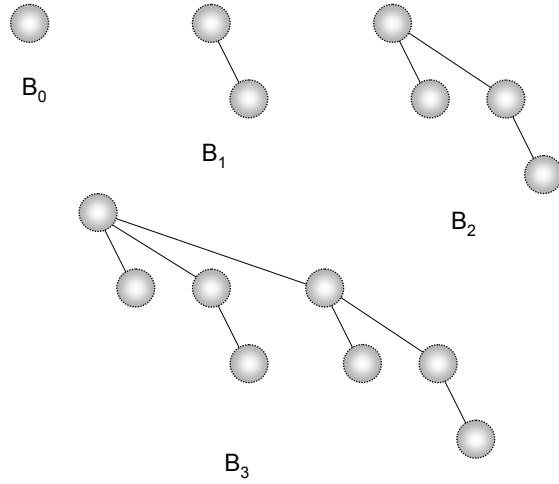
Bilkent University

Computer Engineering Department

Binomial Queues

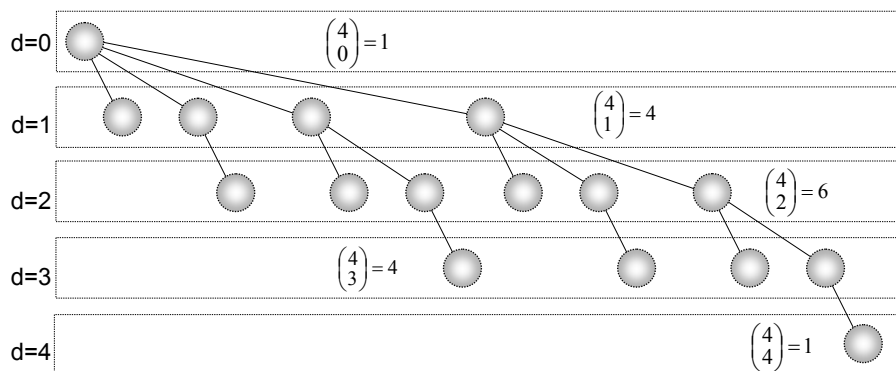
- Binomial Queue
 - A collection of heap-ordered trees (a forest).
 - A tree in the forest is a *binomial tree*.
 - There is *at most one* binomial tree of *every height* in the forest.
- A binomial tree of height 0 has one node.
- A binomial tree of height k has 2^k nodes.
- A binomial tree of height k , B_k , is formed by attaching a binomial tree of height $k-1$, B_{k-1} , to another binomial tree of height $k-1$, B_{k-1} .

Binomial Trees



- A binomial tree B_k consists of root with children B_0, B_1, \dots, B_{k-1}
- Binomial trees of height k have exactly 2^k nodes.
- The number of nodes at depth d is binomial coefficient $\binom{k}{d}$

$k = 4$

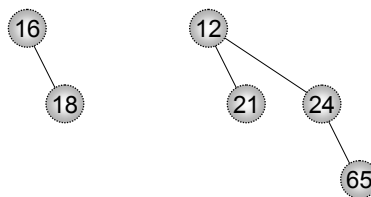


B_4

Heaps and binomial trees

- Assume all binary trees in a forest are in heap order.
- Assume we have at most one binomial tree of any height.
- Then we can represent a heap of any size uniquely by a binomial tree forest.
- Example: heap of size 13 could be represented by forest: B_3, B_2, B_0 ,
- We call this heap ordered binomial forest a **binomial queue**.

Example



H_1

A binomial queue H_1 of size 6 is shown above

Can be represented as: 110

Binomial Queue Operations

- FindMin
- Merge
- Insert
- DeleteMin

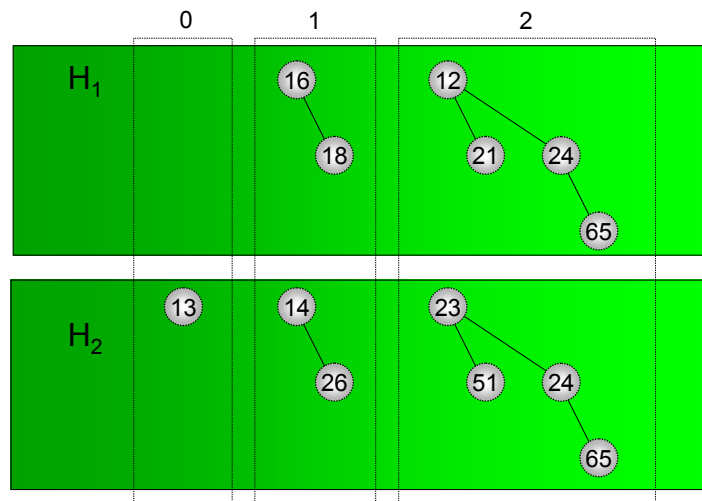
FindMin

- The minimum element in the binomial queue can be found by scanning the roots of all trees in the forest.
- There are at most $\log N$ different trees
- The cost of findMin is therefore $O(\log N)$ in the worst case.

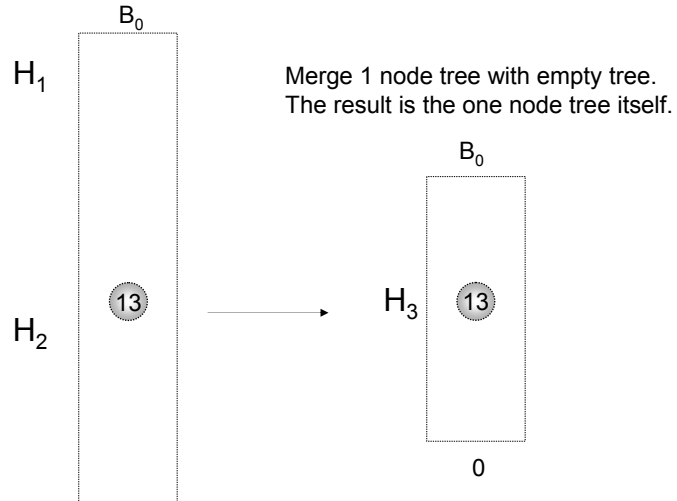
Merge

- Merging two binomial queues is conceptually simple.
- Merge H_1 and H_2
- Just add them in binary.
- Assume H_1 has 6 nodes. H_2 has 7 nodes.
 - $H_3 = \text{node}(H_1) + \text{nodes}(H_2) = 13$.
 - H_1 : 0110
 - H_2 : 0111
 - H_3 : 1101 (this is the resulting bin. heap).

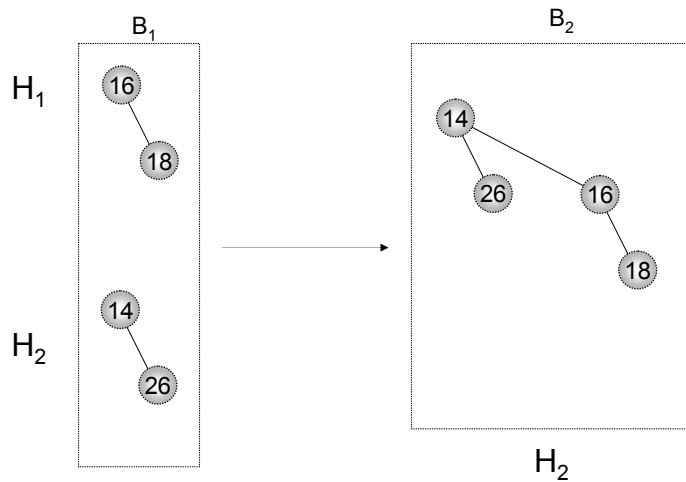
Merge - example

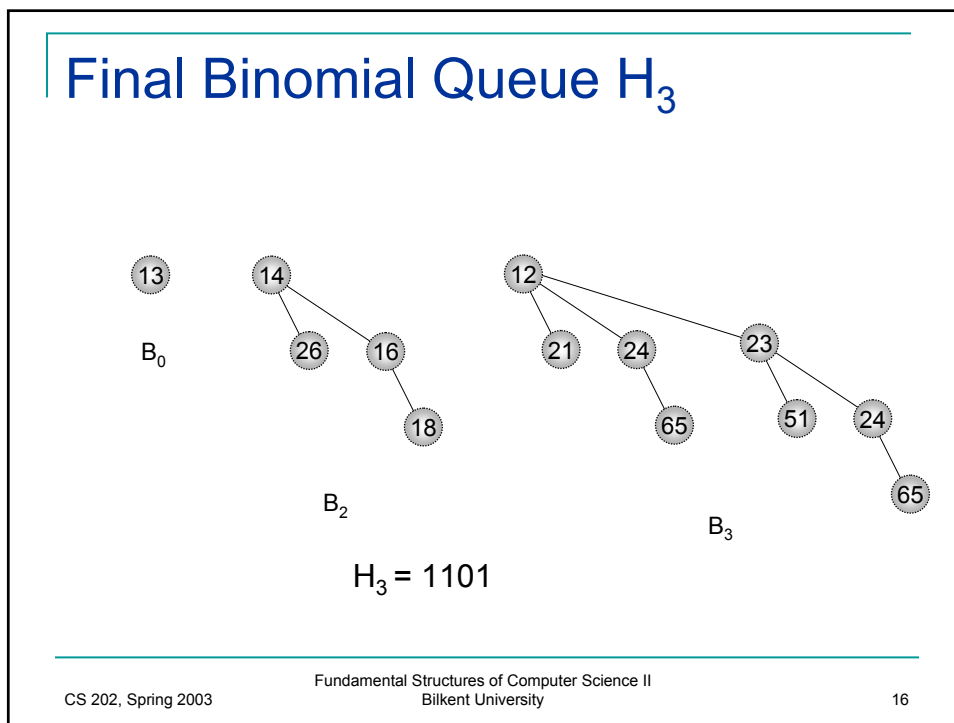
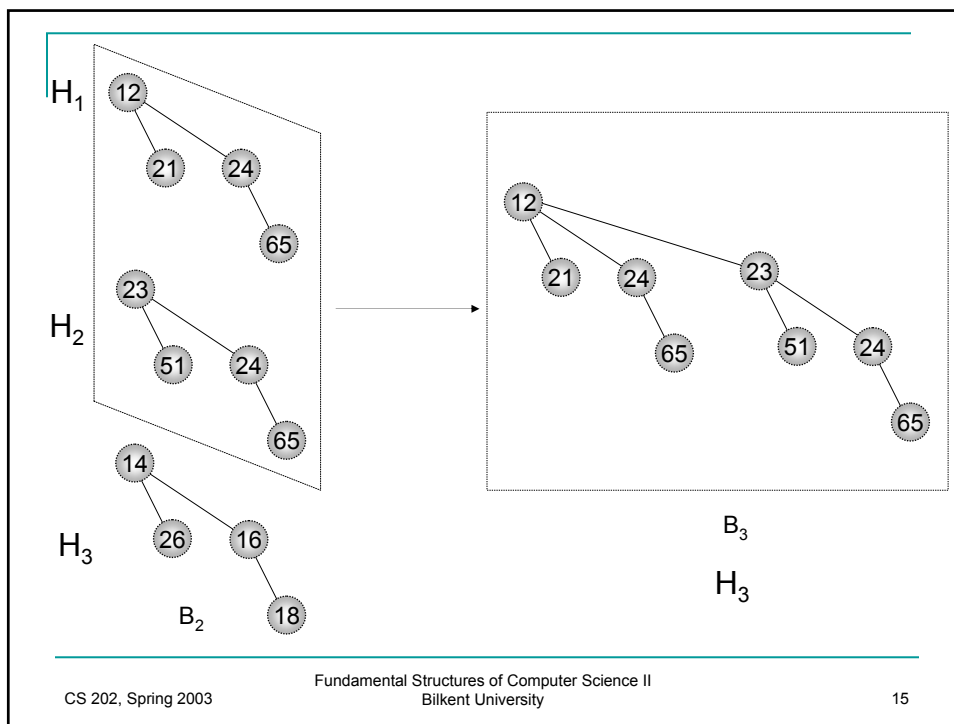


Merge - example - steps



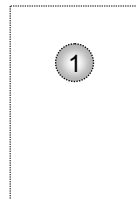
Merge - example - steps



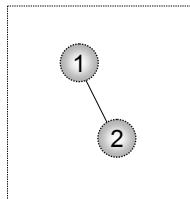


Insert

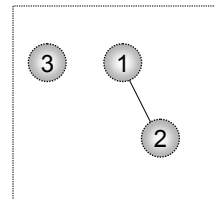
- Insertion is a special case of merge
- Insert an item to a binomial heap H_1
- Make a new heap, H_2 , of one node (item to be inserted)
- Merge H_1 and H_2 .
- Example
 - Insert items 1, 2, 3, 4, 5, 6, 7, in the given order into an empty binomial heap.



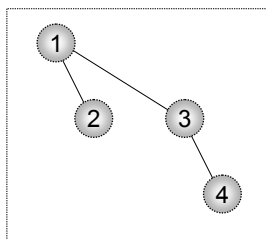
After 1 is inserted



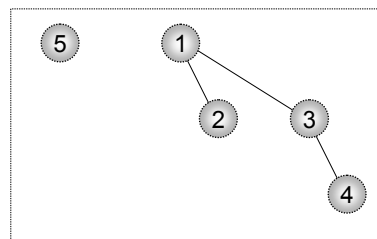
After 2 is inserted



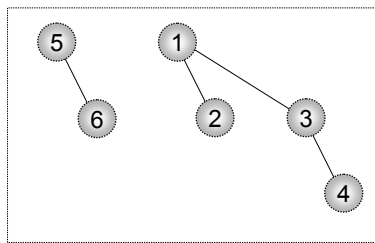
After 3 is inserted



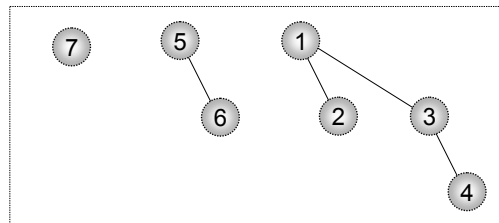
After 4 is inserted



After 5 is inserted



After 6 is inserted

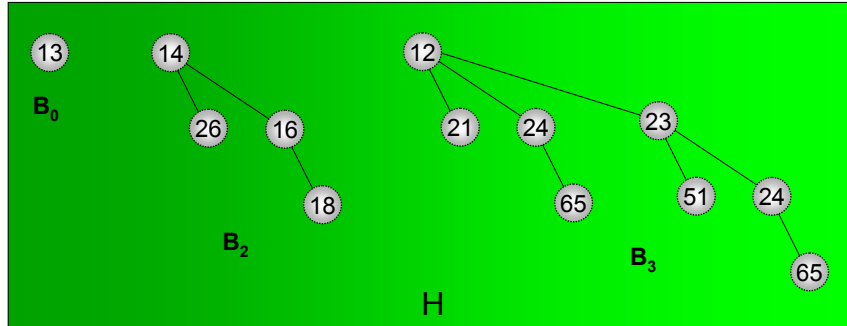


After 7 is inserted

DeleteMin

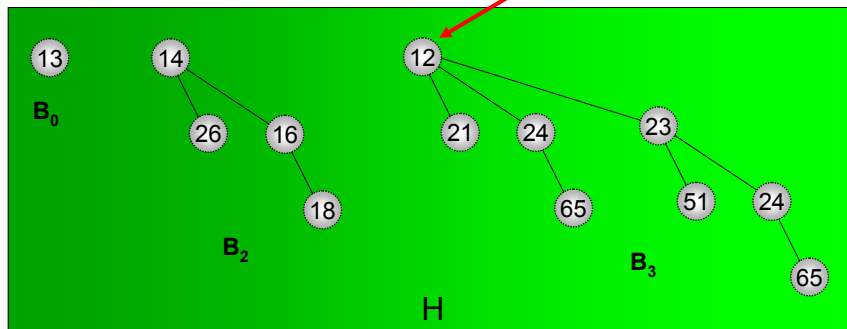
- Sketch of the Algorithm
- Assume we want to delete minimum item from binomial queue H .
 - Find the binomial tree in H that has the minimum root. Let say this is B_k in H .
 - Take tree B_k from H . Let the remaining trees in H make a new binomial queue H_1
 - Remove the root from B_k (root is the minimum that you should return as a result of algorithm).
 - The children of root of B_k make a new heap H_2 that can consists of tree B_0 through B_{k-1}
 - Merge H_1 and H_2 .

DeletMin - Example



Delete the minimum item from the binomial queue above.

DeletMin - Example



The root that has the minimum item is 12 and belongs to tree B_3

