

# Stream Ciphers

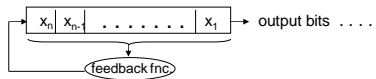
BİL 448/548  
Internet Security Protocols  
Ali Aydın Selçuk

# Stream Ciphers

- Generate a *pseudo-random* key stream & xor to the plaintext.
- Key: The seed of the PRNG
- Traditional PRNGs (e.g. those used for simulations) are not secure.  
E.g., the linear congruential generator:  
$$X_i = a X_{i-1} + b \pmod m$$
for some fixed  $a, b, m$ .
- It passes the randomness tests, but it is predictable if previous output bytes are known.

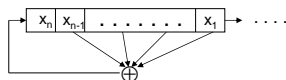
# Linear Feedback Shift Registers

- Feedback shift register:



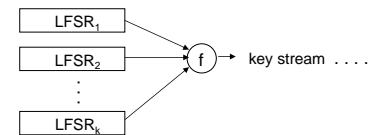
(“register”, “feedback”, “shift”)

- LFSR: Feedback fnc. is linear over  $Z_2$  (i.e., an xor):



- Very compact & efficient in hardware (e.g., SIM cards)

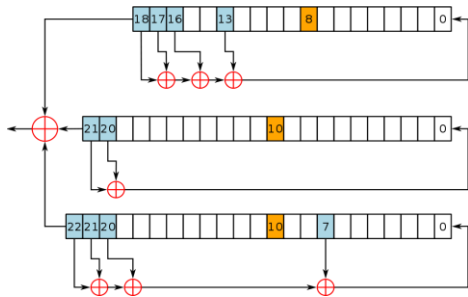
# Stream Ciphers from LFSRs



Desirable properties of  $f$ :

- high non-linearity
- long “cycle period” ( $\sim 2^{n_1+n_2+\dots+n_k}$ )
- low correlation with the input bits

## GSM A5/1



- The A5/1 stream cipher uses three LFSRs.
- A register is clocked if its clocking bit (orange) agrees with one or both of the clocking bits of the other two registers. (majority match)

BI448, A.Selçuk

Stream Ciphers

5

## Software-Oriented Stream Ciphers

- LFSRs are slow in software
- Alternatives:
  - Block ciphers (or hash functions) in CFB, OFB, CTR modes.
  - Stream ciphers designed for software: RC4, SEAL, SALSA20, SOSEMANUK...

BI448, A.Selçuk

Stream Ciphers

6

## RC4 (Rivest, 1987)

- Simple, byte-oriented, fast in s/w.
- Popular: Google, MS-Windows, Apple, Oracle Secure SQL, WEP, WPA, etc.

Algorithm:

- Works on  $n$ -bit words. (typically,  $n = 8$ )
- State of the cipher: A permutation of  $\{0, 1, \dots, N-1\}$ , where  $N = 2^n$ , stored at  $S[0, 1, \dots, N-1]$ .
- Key schedule: Expands the  $\ell$ -byte key (typically 40-256 bits) into the initial state table  $S$ .

BI448, A.Selçuk

Stream Ciphers

7

## RC4 Key Schedule

The key schedule (i.e., initialization) algorithm:

```
// typically  $n = 8$ ,  $\ell = 16$ 
for  $i = 1$  to  $2^n - 1$  do:
     $S[i] \leftarrow i$ 
 $i \leftarrow 0$ ,  $j \leftarrow 0$ 
for  $i = 1$  to  $2^n - 1$  do: {
     $j \leftarrow j + S[i] + K[i \bmod \ell]$ 
     $S[i] \leftrightarrow S[j]$ 
}
```

BI448, A.Selçuk

Stream Ciphers

8

## RC4 Encryption

The encryption (i.e., the PRNG) algorithm:

```
i ← 0, j ← 0
loop: {
  i ← i + 1
  j ← j + S[i]
  S[i] ↔ S[j]
  output S[S[i] + S[j]]
}
```

## IV for Stream Ciphers

- Use of an initialization vector is crucial in a stream cipher.
- Otherwise, the same stream will be produced each time the key is used (i.e., for each packet).
- The cipher may specify how to incorporate the IV. e.g., A5/1 mixes 22-bit frame no. into registers.
- Otherwise, ad hoc methods are used. e.g., WEP uses RC4 with 128-bit  $K' = (IV || K)$  for a 24-bit IV and a 104-bit K.

## Speed of Software Stream Ciphers

(Crypto++ 5.6 benchmarks, 2.2 GHz AMD Opteron 8354.  
March 2009.)

| Algorithm     | Speed (MiByte/s.) |
|---------------|-------------------|
| 3DES / CTR    | 17                |
| AES-128 / CBC | 148               |
| AES-128 / CTR | 198               |
| RC4           | 124               |
| SEAL          | 447               |
| SOSEMANUK     | 767               |
| SALSA20       | 953               |