

Final Exam

January 4, 2010

Question 1. (*60 pts.*) Answer briefly each of the following questions: B

- a. What is the major limitation of a traditional substitution cipher? How do modern block ciphers address it?
- b. Consider the RSA algorithm where the modulus n is a large prime rather than a composite number. Would the encryption be secure? Why/why not?
- c. What is the basic idea of ID-based encryption? Name an advantage and a disadvantage of ID-based encryption against traditional, certificate-based solutions.
- d. What is function sharing? What is its advantage over simple secret sharing?
- e. What is the salt in a password-based authentication system? How does it help in defending against offline dictionary attacks?
- f. Describe how the Kerberos login protocol works. Is it secure against offline password guessing with passive eavesdropping?
- g. What is the advantage of a puzzle scheme for DoS protection over a cookie scheme?
- h. What is the main advantage of EKE-type strong password protocols?
- i. What is an SA in IPsec? Describe how the sender and the receiver finds the SA to be used to process an IP packet.
- j. Describe how 3D-Secure is related to and differs from its predecessor SET.
- k. What is the advantage of using 3D-Secure in on-line purchases?
- l. What are the relative advantages and disadvantages of S/MIME and PGP? What are the environments that would favor each?

Question 2. (20 pts.) Explain the following points regarding Tuomas Aura's talk on Mobile IPv6 security:

- a. What is the problem with the binding updates when no authentication is present?
- b. What is the approach of Aura et al. to this problem?
- c. What is the basic solution they propose?
- d. What is the problem with this solution when satellite links are common?
- e. What would be a simple solution to prevent such passive attacks? (without any major changes to the basic setting such as adding a PKI)

Question 3. (20 pts.) A protocol to establish a fresh session key using long-term, certified Diffie-Hellman public keys is the protocol of Yacobi and Shmueli. The protocol, in a slightly modified form, is as follows:

- The system has a common prime modulus p and a generator g . Each party i has a long-term private key $\alpha_i \in \mathbb{Z}_{p-1}$ and a public key $P_i = g^{\alpha_i} \bmod p$.
- To establish a session key between i and j , party i generates a random $R_i \in \mathbb{Z}_{p-1}$, computes $X_i = \alpha_i + R_i \bmod p - 1$, and sends X_i to j . Similarly, j computes a random $R_j \in \mathbb{Z}_{p-1}$, $X_j = \alpha_j + R_j \bmod p - 1$, and sends X_j to i .
- i computes the session key as

$$K_{i,j} = (g^{X_j} P_j^{-1})^{R_i} \bmod p$$

and j computes

$$K_{j,i} = (g^{X_i} P_i^{-1})^{R_j} \bmod p.$$

- a. Show that the protocol is correct (i.e., $K_{i,j} = K_{j,i}$).
- b. Show that a passive attacker Trudy who has broken a session key $K_{A,B}$ between Alice and Bob can compute any future session keys between these two parties.
- c. Describe a simple addition to the session key computation which will preclude this and any similar attacks on this protocol.

Good luck